COMPARING HIERARCHIES OF TOTAL FUNCTIONALS

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ABSTRACT. In this paper, we will address a problem raised by Bauer, Escardó and Simpson. We define two hierarchies of total, continuous functionals over the reals based on domain theory, one based on an "extensional" representation of the reals and the other on an "intensional" representation. The problem is if these two hierarchies coincide. We will show that this coincidence problem is equivalent to the statement that the topology on the Kleene-Kreisel continuous functionals of a fixed type induced by all continuous functions into the reals is zero-dimensional for each type.

As a tool of independent interest, we will construct topological embeddings of the Kleene-Kreisel functionals into both the extensional and the intensional hierarchy at each type. The embeddings will be hierarchy embeddings as well in the sense that they are the inclusion maps at type 0 and respect application at higher types.

1. INTRODUCTION

There are essentially two ways (with many dialects) to represent the reals as dataobjects using domains. One way is to use approximations to the reals in such a way that when two objects approximate the same real, they are consistent in the domain-theoretical sense. This is called *an extensional approach*. Another way is to represent a real via a sequence of integers, representing some approximating sequence. Two sequences may represent the same real, as in $1.000\cdots$ and $0.999\cdots$, but as data-objects they are quite different and will be considered as inconsistent pieces of information. This is called *the intensional approach*.

We will consider one example of each kind.

Our example of an extensional representation will be via the algebraic domain of closed rational intervals where the reals are represented by ideals of intervals such that the intersection of the sets in the ideal consists of one point.

Our example of an intensional representation will be the binary negative digit representation, essentially representing a real as an infinite sum

$$a + \sum_{i=1}^{\infty} b_i \cdot 2^{-i}$$

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where $a \in \mathbb{Z}$ and each b_i is in $\{-1, 0, 1\}$. This will essentially be an admissible representation of the reals as defined by Weihrauch[23].

With the intensional representation we may consider the representatives for the reals as the total elements of an algebraic domain in a natural way. Thus in both the extensional and the intensional case we may first construct the typed hierarchy of continuous functionals of finite types in the category of algebraic domains, then isolate the hereditarily total functionals in both hierarchies and finally consider the extensional collapse of both hierarchies, see Sections 2 and 4 for details.

The problem is if these two hierarchies coincide.

Why is this an interesting problem?

One of the main motivations for considering typed hierarchies of domains in the first place is to use them for denotational interpretations of programs in some extension of typed λ -calculus. When we add a base type representing the reals, it is because we want to consider programs where reals, or representations of reals, are accepted as data-objects. If the language at hand accepts the reals themselves as data-objects we use an extensional hierarchy, while if it accepts representations for reals, e.g. in the form of data-streams, we use an intensional hierarchy. Typically, *RealPCF* (Escardó [4]) is naturally interpreted over a typed hierarchy of continuous functionals based on the extensional approach while e.g. the approach to typed computability over the reals taken by Di Gianantonio [5, 6, 7], and also by Simpson [21], makes use of an intensional approach.

Since non-termination is an important aspect of computations, it is natural to use algebraic or continuous domains when constructing these hierarchies. As pointed out in e.g. Plotkin [19], an important class of programs are those that terminate on every relevant input, and thus the hereditarily total objects in a typed hierarchy will be of a special interest. Our question is in essence if the notion of a hereditarily total and continuous functional of finite type over the reals is an absolute notion, or if it is a notion that depends on our choice of representations of real numbers.

The precise version of the problem we address was first formulated by Bauer, Escardó and Simpson [3]. They proved that at the first three levels (types 0 - 2), the hierarchies coincide. Moreover, they showed that if the topology on the Kleene-Kreisel continuous functionals of type 2 is zero-dimensional, then the hierarchies also coincide for types at level 3. They further showed that the assumption of zero-dimensionality may be weakened, but that equality of the extensional and intensional hierarchies for types at level 3 will have consequences for the topology of $Ct_{\mathbb{N}}(2)$.

Since then, Normann [17] analyzed the intensional hierarchy further, and described a representation of it via domains such that the total objects of each type form dense subsets of the underlying domains.

Both the hierarchy of continuous total functionals based on the extensional representation of the reals and the one based on the intensional representation can be viewed as natural analogues of the Kleene-Kreisel continuous functionals [11, 12], where \mathbb{N} is replaced by \mathbb{R} . Our first result, which is of independent interest, is that the typed structure of the Kleene-Kreisel continuous functionals may be continuously embedded into the extensional hierarchy over the reals. We will also prove a similar result for the intensional hierarchy. These results will be stated more precisely in due course. In this paper, these embeddings will be used as tools in analyzing the coincidence problem. The typed hierarchies in question are obtained as quotient spaces of hereditarily total objects under the relation of being equivalent. In the extensional case, equivalence will be the same as consistency for hereditarily total objects.

We will make use of the approach from Normann [17]. One consequence of the density theorem in [17] is that consistency will be an equivalence relation on the hereditarily total objects, and that consistent total objects will be equivalent in the sense of representing the same functional. We will study the intermediate hierarchy of quotient spaces dividing the hereditarily total objects of our alternative hierarchy just by consistency instead of full equivalence, and see that the topology of these intermediate spaces share the relevant properties of the Kleene-Kreisel functionals, via mutual topological embeddings. This intermediate hierarchy, the embedding results, the density theorems referred to above and an approximation lemma proved in Normann [18] are used together with a technique from Bauer, Escardó and Simpson [3] in order to link the coincidence problem to a problem about the topology of the Kleene-Kreisel functionals. This will be made more precise later.

Organization. In Section 2 we will give the construction of the Kleene-Kreisel continuous functionals, the $Ct_{\mathbb{R}}(k)$ hierarchy, and of the analogue hierarchy of extensional functionals over the reals, the $Ct_{\mathbb{R}}^{E}(k)$ -hierarchy. We will also state, and to some extent prove, the relevant properties of these hierarchies and results from topology in general that we will need. There is hardly any original material in this section.

In Section 3 we will state and prove the embeddability of the $Ct_{\mathbb{N}}(k)$ -hierarchy into the $Ct_{\mathbb{R}}^{E}$ -hierarchy.

In Section 4 we will introduce the intensional $Ct^{I}_{\mathbb{R}}$ -hierarchy and the *smoothened version*, the equivalent $Ct^{S}_{\mathbb{R}}$ -hierarchy, and we will prove a conditional coincidence theorem.

In Section 5 we give a full characterization of the coincidence problem by proving the converse of the main theorem in Section 4.

In Section 6 we discuss some further problems, and in the appendix (Section 7) we will prove a special case of the approximation lemma from Normann [18].

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2. Background

We will assume that the reader is familiar with the theory of Scott Domains, or bounded complete algebraic domains. In this paper all domains in question will be algebraic and bounded complete, i.e. each bounded set will have a least upper bound. These properties will not necessarily be repeated when assumed. We recommend Stoltenberg-Hansen & al. [22], Abramsky and Jung [1] or Gierz & al. [8] for an introduction to domain theory. For an introduction to the domain theoretical approach to the Kleene-Kreisel continuous functionals [11, 12] we suggest the handbook paper Normann [15].

We will describe the construction of the Kleene-Kreisel continuous functionals and the corresponding hierarchy over the reals based on the extensional representation mainly by setting the notation to be used in the paper.

We will restrict our attention to the pure types. In this paper we will let \mathbb{N} denote the non-negative integers.

Definition 2.1. Let $N(0) = \mathbb{N}_{\perp}$, with $\overline{N}(0) = \mathbb{N}$ and $n \approx_0^N m$ if and only if $n = m \in \mathbb{N}$. Let $N(k+1) = N(k) \to N(0)$ in the category of algebraic domains. For $f, g \in N(k+1)$, let

$$f \approx_{k+1}^{N} g \Leftrightarrow \forall a \in N(k) \forall b \in N(k) (a \approx_{k}^{N} b \Rightarrow f(a) \approx_{0}^{N} f(b))$$

Let $\bar{N}(k+1) = \{ f \in N(k+1) ; f \approx_{k+1}^{N} f \}.$

 \approx_k^N will be a partial equivalence relation. The elements of \bar{N}_k will be called the *hered-itarily total functionals*, and equivalence will mean that equivalent functionals will give us the same well defined output to a hereditarily total input.

Since \approx_k^N is both symmetric and transitive, we will have that if x and y are elements of N(k) such that $x \approx_k^N y$, then $x \in \overline{N}(k)$.

Proposition 2.2. For
$$x, y \in N(k)$$
, we have that $x \approx_k^N y \Leftrightarrow x \sqcap y \in N(k)$.
For $x, y \in \overline{N}(k)$ we have that $x \approx_k^N y \Leftrightarrow x$ and y are consistent.

For a proof, see e.g. Normann [15]. The first part of this proposition was originally proved in Longo and Moggi [14], and the second part is a consequence of the domain-theoretical version of the Kleene-Kreisel Density Theorem, see Proposition 2.7.

Definition 2.3. By recursion on k we define the set $Ct_{\mathbb{N}}(k)$ and the projection map ρ_k^N : $\overline{N}(k) \to Ct_{\mathbb{N}}(k)$ as follows:

 $\rho_0^N(n) = n, Ct_{\mathbb{N}}(0) = \mathbb{N}.$

As an induction hypothesis, an arbitrary element of $Ct_{\mathbb{N}}(k)$ will be of the form $\rho_k^N(x)$ where $x \in \overline{N}(k)$. If $f \in \overline{N}(k+1)$, we let

$$\rho_{k+1}^N(f)(\rho_k^N(x)) = f(x).$$

This is well defined by the definition of \approx_{k+1}^{N} , assuming that ρ_{k}^{N} identifies exactly \approx_{k}^{N} -equivalent objects, and then ρ_{k+1}^{N} will identify exactly \approx_{k+1}^{N} -equivalent objects. Let $Ct_{\mathbb{N}}(k+1) = \{\rho_{k+1}^{N}(f); f \in \overline{N}(k+1)\}.$

 $Ct_{\mathbb{N}}(k)$ is known as the Kleene-Kreisel continuous functionals of type k.

The topology on $Ct_{\mathbb{N}}(k)$ will be the finest topology such that $\rho_k^{\mathbb{N}}$ is continuous. Then $Ct_{\mathbb{N}}(k+1)$ will consist of exactly all continuous maps $F: Ct_{\mathbb{N}}(k) \to \mathbb{N}$.

Using similar constructions, we will now define a hierarchy of functionals over the reals. This will be based on the extensional representation of reals, and we will use the letter E for *extensional* to denote this hierarchy.

Definition 2.4. Let E(0) be the algebraic domain of ideals over

$$\{\mathbb{R}\} \cup \{[p,q] \; ; \; p \in \mathbb{Q} \land q \in \mathbb{Q} \land p \le q\}$$

where the intervals are ordered by reverse inclusion.

Let $E(k+1) = E(k) \rightarrow E(0)$.

If α is an ideal in E(0), then $\cap \alpha \neq \emptyset$. We let

 $\alpha \approx_0^E \beta \Leftrightarrow \cap \alpha = \cap \beta = \{x\} \text{ for some } x \in \mathbb{R}.$

We define \approx_k^E by recursion on k in analogy with the definition of \approx_k^N , and let $\bar{E}(k) = \{x \in E(k) ; x \approx_k^E x\}$.

We then have

Proposition 2.5. (Normann [16])

For $x, y \in E(k)$ we have that $x \approx_k^E y \Leftrightarrow x \sqcap y \in \overline{E}(k)$. For $x, y \in \overline{E}(k)$ we have that $x \approx_k^E y \Leftrightarrow x$ and y are consistent.

Definition 2.6. We define $Ct^E_{\mathbb{R}}(k)$ and $\rho^E_k : \overline{E}(k) \to Ct^E_{\mathbb{R}}$ in analogy with Definition 2.3 as follows: We let $\rho^E_0(\alpha) = x$ if $\{x\} = \cap \alpha$.

follows: We let $\rho_0^E(\alpha) = x$ if $\{x\} = \cap \alpha$. We let $\rho_{k+1}^E(x)(\rho_k^E(y)) = \rho_0^E((x)(y))$. We let $Ct_{\mathbb{R}}^E(k) = \{\rho_k^E(x) \ ; \ x \in \bar{E}(k)\}.$

The domains E(k) are special instances of domains $E(\sigma)$ for all types σ . Then the *E*-hierarchy may be used to implement Escardó's *RealPCF* [4], though the approach via continuous domains is the one used originally.

We define the topology on $Ct^{E}_{\mathbb{R}}(k)$ as the finest topology making $\rho_{E}(k)$ continuous. Then $Ct^{E}_{\mathbb{R}}(k+1)$ is exactly the set of continuous functions from $Ct_{\mathbb{R}}(k)$ to \mathbb{R} .

Proposition 2.7.

- a) For each $k \in \mathbb{N}$, $\overline{N}(k)$ is dense in N(k).
- b) For each $k \in \mathbb{N}$, E(k) is dense in E(k).

Part a) is the domain-theoretical version of the Kleene-Kreisel Density Theorem. Part b) is proved in [16]

In this paper we will work with ω - algebraic domains A (meaning that the set of compacts is countable), a set $\bar{A} \subseteq A$ of "total" objects, an equivalence relation \approx_A on \bar{A} and (essentially) the quotient topology on $\bar{A} \approx_A$.

This topology will be *sequential*, which means that it is generated from the set of convergent sequences with limits.

Since A is ω -algebraic, we see that any open covering of a subset of A/\approx_A can be replaced by a countable sub-covering, i.e. the spaces are *hereditarily Lindelöf*.

These are facts of basic general topology.

In the sequel we will also make use of the following fact:

Lemma 2.8. Let X be a topological space, $\rho : X \to Y$ an onto map and let Y be equipped with the identification topology, i.e. $O \subseteq Y$ is open if and only if $\rho^{-1}[O]$ is open in X.

Let $Z \subseteq Y$ be closed, let $Z_1 = \rho^{-1}[Z]$, and let ρ_Z be ρ restricted to Z_1 .

Then the identification topology on Z induced from ρ_Z and the subspace topology on Z_1 will coincide with the subspace topology on Z.

The proof is elementary and is left for the reader.

Following Scarpellini [20] on one hand (see also [10] or [15]) and Normann [16] on the other, we also have that the hierarchies $\{Ct_{\mathbb{N}}(k)\}_{k\in\mathbb{N}}$ and $\{Ct_{\mathbb{R}}^{E}(k)\}_{k\in\mathbb{N}}$ can be defined in the category of Kuratowski limit spaces, see Kuratowski [13]. We have the canonical limit structures on \mathbb{N} and \mathbb{R} resp. Then, by recursion on k we may define $Ct_{\mathbb{N}}(k)$ and $Ct_{\mathbb{R}}^{E}(k)$ with limit structures as follows, where Ct(k) may stand for both sets:

 $F \in Ct(k+1)$ if $F : Ct(k) \to Ct(0)$ and for all $\{a_n\}_{n \in \mathbb{N}}$ and a from Ct(k), $a = \lim_{n \to \infty} a_n \Rightarrow F(a) = \lim_{n \to \infty} F(a_n)$, i.e., F is sequence continuous.

 $F = \lim_{n \to \infty} F_n$ if for all $\{a_n\}_{n \in \mathbb{N}}$ and a from Ct(k), $a = \lim_{n \to \infty} a_n \Rightarrow F(a) = \lim_{n \to \infty} F_n(a_n)$.

For both hierarchies of quotient spaces it is also the case that the convergent sequences of quotients with limits are exactly the sequences obtained by taking a convergent sequence at the domain level and then take the quotients. For the Kleene-Kreisel continuous functionals there is a third characterization of the convergent sequences as well. The following proposition summarizes all this. Part a) was proved in Hyland [10], for an exposition see [15]. Part b) was proved in Normann [16].

Proposition 2.9.

- a) Let $k \in \mathbb{N}$, $\{f_i\}_{i \in \mathbb{N}}$ be a sequence from $Ct_{\mathbb{N}}(k+1)$ and $f \in Ct_{\mathbb{N}}(k+1)$. Then the following are equivalent:
 - i) $f = \lim_{i \to \infty} f_i$ in the topology on $Ct_{\mathbb{N}}(k+1)$.
 - ii) Whenever $a = \lim_{i \to \infty} a_i$ in $Ct_{\mathbb{N}}(k)$, then $f(a) = \lim_{i \to \infty} f_i(a_i)$ in \mathbb{N} .
 - iii) There is a convergent sequence with limit $g = \lim_{i \to \infty} g_i$ from $\bar{N}(k+1)$ such that $f = \rho_{k+1}^N(g)$ and $f_i = \rho_{k+1}^N(g_i)$ for each $i \in \mathbb{N}$
 - vi) There is a continuous modulus g of convergence, i.e. $g \in Ct_{\mathbb{N}}(k+1)$ such that

$$\forall a \in Ct_{\mathbb{N}}(k) \forall i \ge g(a)(f_i(a) = f(a))$$

- b) Let $k \in \mathbb{N}$, $\{f_i\}_{i \in \mathbb{N}}$ be a sequence from $Ct^E_{\mathbb{R}}(k+1)$ and $f \in Ct^E_{\mathbb{R}}(k+1)$. Then the following are equivalent:
 - i) $f = \lim_{i \to \infty} f_i$ in the topology on $Ct^E_{\mathbb{R}}(k+1)$.
 - ii) Whenever $a = \lim_{i \to \infty} a_i$ in $Ct^E_{\mathbb{R}}(k)$, then $f(a) = \lim_{i \to \infty} f_i(a_i)$ in \mathbb{R} .
 - iii) There is a convergent sequence with limit $g = \lim_{i \to \infty} g_i$ from $\overline{E}(k+1)$ such that $f = \rho_{k+1}^E(g)$ and $f_i = \rho_{k+1}^E(g_i)$ for each $i \in \mathbb{N}$.

We will not prove these results in detail, but, as we will see, they follow from the next proposition and a similar proposition for the reals:

Proposition 2.10. Let (X, \sqsubseteq) be an ω -algebraic domain and let $\overline{X} \subseteq X$ be a subset that is uppwards closed. If $f: \overline{X} \to \mathbb{N}$ is continuous in the subspace topology on \overline{X} , then f can be extended to a continuous $g: X \to \mathbb{N}_{\perp}$.

Proposition 2.10 is due to the author, but the proof was not published. If we replace \mathbb{N} by \mathbb{R} , the proposition is proved as Lemma 6.2 in Normann [16], and the proof from [16] works in this simpler case as well. In the present paper, the method of proof is employed in the proof of Theorem 7.1, in defining Y_n^f from X_n^f . The argument is also used in Normann [18].

Corollary 2.11. In Proposition 2.10 we may replace \mathbb{N} with $\mathbb{N}^{\mathbb{N}}$ and \mathbb{N}_{\perp} with $\mathbb{N}_{\perp} \to \mathbb{N}_{\perp}$.

Proof. Let $g: \overline{X} \to \mathbb{N}^{\mathbb{N}}$ be continuous.

Let $g_1(x,n) = g(x)(n)$ and apply Proposition 2.10 to $g_1: \overline{X} \times \mathbb{N} \to \mathbb{N}$.

In this paper, function spaces, and subspaces of function spaces, will play an important part. There is no canonical topology on a function space, so we will now discuss the topology we will use for the relevant cases.

For the purpose of this discussion, let (X, \bar{X}, \approx_X) be an ω -algebraic domain X with a totality \bar{X} satisfying the requirement of Proposition 2.10 and let \approx_X be an equivalence relation on \bar{X} . Let $\mathcal{T}(X)$ be \bar{X}/\approx_X with the quotient topology. Let $\mathcal{T}(X) \to \mathbb{N}$ be the set of continuous functions from $\mathcal{T}(X)$ to \mathbb{N} . We will define a "default" topology on $\mathcal{T}(X) \to \mathbb{N}$ and establish some properties of this topology. These properties will extend to $\mathcal{T}(X) \to \mathbb{N}^{\mathbb{N}}$, and then to $\mathcal{T}(X) \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$, for which we will use them. (\mathbb{N}^+ is the set of positive integers.)

Let $f \in \mathcal{T}(X) \to \mathbb{N}$. By Proposition 2.10 there is a continuous $g: X \to \mathbb{N}_{\perp}$ such that for each $a \in \mathcal{T}(X)$ and each $x \in a$ we have that g(x) = f(a). Let $Y = X \to \mathbb{N}_{\perp}$ in the category of algebraic domains. Let $g \in Y$ and $h \in Y$. Let

$$g \approx_Y h \Leftrightarrow \forall x \in \bar{X} \forall y \in \bar{X} (x \approx_X y \to g(x) = h(y) \in \mathbb{N}).$$

Then \approx_Y will be a partial equivalence relation, and we let

$$\bar{Y} = \{g \in Y \; ; \; g \approx_Y g\}.$$

We use the quotient topology on \overline{Y} / \approx_Y and the canonical 1-1 correspondence between \overline{Y} / \approx_Y and $\mathcal{T}(X) \to \mathbb{N}$ to define the topology on $\mathcal{T}(X) \to \mathbb{N}$.

Definition 2.12. Let $(X, \overline{X}, \approx_X)$ be as above.

We say that $\mathcal{T}(X)$ accepts liftings of convergent sequences if whenever $a = \lim_{n \to \infty} a_n$ in $\mathcal{T}(X)$ then there is an $x \in a$ and an $x_n \in a_n$ for each n such that $x = \lim_{n \to \infty} x_n$.

Lemma 2.13. Let $(X, \overline{X}, \approx_X)$ and $\mathcal{T}(X)$ be as above and assume that $\mathcal{T}(X)$ accepts liftings of convergent sequences. Let $(Y, \overline{Y}, \approx_Y)$ and $\mathcal{T}(X) \to \mathbb{N}$ be as constructed.

- a) $\mathcal{T}(Y)$ accepts liftings of convergent sequences.
- b) If $f \in \mathcal{T}(X) \to \mathbb{N}$ and $f_n \in \mathcal{T}(X) \to \mathbb{N}$ for each $n \in \mathbb{N}$, then the following are equivalent:
 - i) $f = \lim_{n \to \infty} f_n$.
 - ii) Whenever $a = \lim_{n \to \infty} a_n$ in $\mathcal{T}(X)$, then $f(a) = \lim_{n \to \infty} f_n(a_n)$.

Proof. Let Ω be the domain with compacts \bot and n and n^* for $n \in \mathbb{N}$ (n^* is just a formal object), where $\bot \sqsubseteq a$ for all compacts $a, n \sqsubseteq m$ if and only if $n = m, n^* \sqsubseteq m^*$ if and only if $n \leq m$ and $n^* \sqsubset k$ if and only if $n \leq k$. Ω is known as the "lazy natural numbers".

We will let ω be the maximal ideal generated from $\{n^*; n \in \mathbb{N}\}$, and notation-wise we will also use "n" for the ideal generated by $\{n\}$. Then $\omega = \lim_{n \to \infty} n$. Let $\overline{\Omega} = \mathbb{N} \cup \{\omega\}$.

The point is that a convergent sequence $f_{\omega} = \lim_{n \to \infty} f_n$ of continuous functions from $\mathcal{T}(X)$ to \mathbb{N} will be the continuous image of $n \mapsto f_n$ where $n \in \overline{\Omega}$.

a) Let $f_{\omega} = \lim_{n \to \infty} f$. Let $F(n, x) = f_{n(x)}$ for $n \in \Omega$.

By Proposition 2.10 there is a continuous $G : \Omega \times X \to \mathbb{N}$ such that $G(n, x) = f_n([x])$ (where [x] is the equivalence class of x) for all $n \in \overline{\Omega}$ and all $x \in \overline{X}$.

Let $g_n(x) = G(n, x)$. Then $g_{\omega} = \lim_{n \to \infty} g_n$. This will be a lifting of the convergent sequence.

b) Let $f_{\omega} = \lim_{n \to \infty} f_n$ and let $a_{\omega} = \lim_{n \to \infty} a_n$.

By the assumption that $\mathcal{T}(X)$ accepts liftings of convergent sequences and by a) of this lemma, it follows that $f(a) = \lim_{n \to \infty} f_n(a_n)$. Assume now that $f_{\omega}(a_{\omega}) = \lim_{n \to \infty} f_n(a_n)$ whenever $a_{\omega} = \lim_{n \to \infty} a_n$.

This means that

$$F(n,a) = f_n(a)$$

is sequence-continuous on $\overline{\Omega} \times \mathcal{T}(X)$ (the sequence topology may be finer than the product topology).

Let Z be the domain $\Omega \times X$, $\overline{Z} = \{(n, x) ; x \in \overline{X} \land n \in \overline{\Omega}\}$ and let $(n, x) \approx_Z (m, y) \Leftrightarrow n = m \land x \approx_X y$. Then the quotient topology on \overline{Z} / \approx_Z is sequential, and is essentially the same as the sequence topology on $\overline{\Omega} \times \mathcal{T}(X)$.

Then, by Proposition 2.10 there is a continuous $G : \Omega \times X \to \mathbb{N}_{\perp}$ such that whenever $a \in \mathcal{T}(X), x \in a$ and $n \in \overline{\Omega}$ we have that

$$G(n,x) = F(n,a) = f_n(a).$$

Let $g_n(x) = G(n,x)$ for $n \in \overline{\Omega}$ and $x \in X$. Then $g_{\omega} = \lim_{n \to \infty} g_n$ and we may conclude that $f_{\omega} = \lim_{n \to \infty} f_n$.

Remark 2.14. As pointed out, a convergent sequence $f_{\omega} = \lim_{n \to \infty} f_n$ corresponds to a continuous map $f : \overline{\Omega} \to (\mathcal{T}(X) \to \mathbb{N})$

By the standard isomorphism, f corresponds to a map $\hat{f}: \bar{\Omega} \times \mathcal{T}(X) \to \mathbb{N}$.

 \hat{f} needs not be continuous in the product topology, but when $\mathcal{T}(X)$ accepts liftings of convergent sequences, \hat{f} is continuous in the quotient topology on $\bar{\Omega} \times \mathcal{T}(X)$. This will also hold when we replace \mathbb{N} by \mathbb{R} .

A set A in a topological space T is called *clopen* if A is both closed and open. We will let $\mathcal{Z}(T)$ be the subtopology generated by the clopen sets. If $f \in Ct_{\mathbb{N}}(k+1)$ for some k, f will be continuous with respect to $\mathcal{Z}(Ct_{\mathbb{N}}(k))$.

Definition 2.15. Let $A \subseteq Ct_{\mathbb{N}}(k)$.

A is a Π_1^0 -set if there are clopen sets A_n such that $A = \bigcap_{n \in \mathbb{N}} A_n$.

In this case the $\tilde{\Pi}_n^0$ -sets are exactly the sets closed in $\mathcal{Z}(Ct_{\mathbb{N}}(k))$. When A is a $\tilde{\Pi}_1^0$ -set in $Ct_{\mathbb{N}}(k)$, then the topology on A induced from the topology on $Ct_{\mathbb{N}}(k)$ will coincide with the quotient topology induced by

$$\bar{A} = \{ x \in \bar{N}(k) \; ; \; \rho_k^{\mathbb{N}}(x) \in A \}$$

and $\approx_k^{\mathbb{N}}$, see Lemma 2.8.

If $A \subseteq Ct^E_{\mathbb{R}}(k)$ is an arbitrary subset, we will consider the quotient topology on A induced by ρ^E_k restricted to $(\rho^E_k)^{-1}(A)$.

 $\{f: A \to \mathbb{R}; f \text{ is continuous}\}$ will have a default topology in analogy with $\mathcal{T}(X) \to \mathbb{N}$, and by the \mathbb{R} -versions of Proposition 2.10 and its consequences, this will be the finest topology where $f = \lim_{n \to \infty} f_n$ exactly when

$$\forall a \in A \forall \{a_n\}_{n \in \mathbb{N}} \in A^{\mathbb{N}}(a = \lim_{n \to \infty} a_n \Rightarrow f(a) = \lim_{n \to \infty} f_n(a_n)).$$

We use this topology in Proposition 2.16 a) below. Part a) is the Approximation Lemma, i.e. Theorem 2, from Normann [18]. Part b) is essentially a special case of a), and will be proved in the Appendix for the sake of completeness.

Proposition 2.16.

- a) Let $A \subseteq Ct^E_{\mathbb{R}}(k)$ and let $f: A \to \mathbb{R}$ be continuous. Then, continuously in f, there are $f_n \in Ct^E_{\mathbb{R}}(k+1)$ such that whenever $x \in A$ and $x = \lim_{n \to \infty} x_n$ with each $x_n \in Ct^E_{\mathbb{R}}(k)$ we have that $f(x) = \lim_{n \to \infty} f_n(x_n)$.
- b) Let $A \subseteq Ct_{\mathbb{N}}(k)$ and let $f : A \to \mathbb{N}$ be continuous. Then, continuously in f, there are $f_n \in Ct_{\mathbb{N}}(k+1)$ such that whenever $x \in A$ and $x = \lim_{n \to \infty} x_n$ with each $x_n \in Ct_{\mathbb{N}}(k)$ we have that $f(x) = \lim_{n \to \infty} f_n(x_n)$.

3. A HIERARCHY OF EMBEDDINGS

3.1. Aim and consequences. N(k) and E(k) are examples of domains with totalities $\overline{N}(k)$ and $\overline{E}(k)$ resp. A continuous map $\pi : N(k) \to E(k)$ is then called *total* if π maps $\overline{N}(k)$ into $\overline{E}(k)$.

We are operating with equivalence relations \approx_k^N and \approx_k^E on $\overline{N}(k)$ and $\overline{E}(k)$ coinciding with consistency, and then a continuous total map $\pi : N(k) \to E(k)$ will induce a continuous map $\overline{\pi} : Ct_{\mathbb{N}}(k) \to Ct_{\mathbb{R}}^E(k)$. In this section we will prove the following

Theorem 3.1. For each $k \in \mathbb{N}$ there is a total, continuous map

$$\pi_k: N(k) \to E(k)$$

such that

- i) $\pi_0: N(0) \to E(0)$ sends a number n to its representative in E(0).
- ii) For each $k \in \mathbb{N}$, for each $f \in \overline{N}(k+1)$ and $a \in \overline{N}(k)$ we have

$$\pi_0(f(a)) = \pi_{k+1}(f)(\pi_k(a)).$$

Before entering the proof of the theorem, we will establish some consequences.

Corollary 3.2. For each $k \in \mathbb{N}$ there is an injective, continuous map

$$\bar{\pi}_k: Ct_{\mathbb{N}}(k) \to Ct_{\mathbb{R}}^E(k)$$

such that

- i) $\bar{\pi}_0$ is the standard inclusion map from \mathbb{N} to \mathbb{R} .
- ii) For each $k \in \mathbb{N}$, each $F \in Ct_{\mathbb{N}}(k+1)$ and each $a \in Ct_{\mathbb{N}}(k)$ we have that $F(a) = Ct_{\mathbb{N}}(k)$ $\bar{\pi}_{k+1}(F)(\bar{\pi}_k(a)).$

Remark 3.3. Independently, Bauer and Simpson [2] gave a proof for Corollary 3.2 for $k \leq 2$. Their result is stronger in the sense that it is proved in intuitionistic logic for constructive analysis.

Another important consequence is the following

Corollary 3.4. Let $\bar{\pi}_k$ be obtained from Theorem 3.1 as in Corollary 3.2.

Then the range of $\bar{\pi}_k$ is a closed subset of $Ct^E_{\mathbb{R}}(k)$ homeomorphic to $Ct_{\mathbb{N}}(k)$.

Proof. By recursion on k we will define continuous, partial inverses $\pi_k^{-1}: E(k) \to N(k)$ as follows:

- If $n \in \mathbb{N}$ and $[p,q] \subseteq (n \frac{1}{3}, n + \frac{1}{3})$, we let $\pi_0^{-1}([p,q]) = n$, while $\pi_0^{-1}([p,q]) = \perp_{\mathbb{N}}$ if the above rule does not apply.
- If $q \in E(k+1)$ and $a \in N(k)$ we let

$$\pi_{k+1}^{-1}(g)(a) = \pi_0^{-1}(g(\pi_k(a))).$$

By induction, i) and ii) below follow from the construction:

- i) If $k \in \mathbb{N}$ and $a \in N(k)$ then $a = \pi_k^{-1}(\pi_k(a))$. ii) If $a \in \bar{E}(k)$, $b \in \bar{E}(k)$, $a \approx_k^E b$ and $\pi_k^{-1}(a) \in \bar{N}(k)$, then $\pi_k^{-1}(b) \approx_k^N \pi_k^{-1}(a)$ and consequently $\pi_k^{-1}(b) \in \bar{N}(k)$.

This shows that $Ct_{\mathbb{N}}(k)$ is homeomorphic to the range of $\bar{\pi}_k$ with the quotient topology. We have to show that this range is closed, and then the rest of the corollary will follow from Lemma 2.8. In order to prove that the range is closed, we use that the topology is sequential, see Proposition 2.9.

The range of $\bar{\pi}_0$ is just the closed subset \mathbb{N} of \mathbb{R} .

If $f \in \overline{E}(k+1)$ we have that $\pi_{k+1}^{-1}(f) \in \overline{N}(k+1)$ if

$$\forall a \in N(k) (f(\pi_k(a)) \in \mathbb{N}).$$

Let $g = \lim_{i \to \infty} g_i$ in $Ct_{\mathbb{R}}(k+1)$ such that each g_i is in the range of $\overline{\pi}_{k+1}$. We will show that the range of $\bar{\pi}_{k+1}$ is closed by showing that g is in the range of $\bar{\pi}_{k+1}$.

By Proposition 2.9, let $h = \lim_{i\to\infty} h_i$ where $h \in \overline{E}(k+1)$, each $h_i \in \overline{E}(k+1)$, $g = \rho_{k+1}^E(h)$ and each $g_i \in \rho_{k+1}^E(h_i)$.

Then $\forall a \in \overline{N}(k)(h(\pi_{k+1}(a)) \in \mathbb{N})$, so $\pi_{k+1}^{-1}(h)$ is defined. Then

$$\pi_{k+1}(\pi_{k+1}^{-1}(h)) = \lim_{i \to \infty} \pi_{k+1}(\pi_{k+1}^{-1}(h_i)) = \lim_{i \to \infty} h_i = h.$$

It follows that h is in the range of π_{k+1} , so g is in the range of $\bar{\pi}_{k+1}$, and the corollary is proved.

During the construction we will observe that π_k is computable such that if $a \in N_k$ is not total, then $\pi_k(a)$ is not total. By Kreisel [12] we know that $\bar{N}(k+1)$ is complete Π_k^1 for $k \ge 1$, see also Normann [15]. It is strait-forward to show by induction on k that $\bar{E}(k+1)$ is a Π_k^1 -set for $k \ge 1$. We then obtain

Corollary 3.5. For $k \ge 1$, $\overline{E}(k+1)$ is complete Π_k^1 .

The following result is proved using Proposition 2.16 and Corollary 3.2. By a suitable adjustment of Proposition 2.16, the use of Corollary 3.2 may be avoided.

Theorem 3.6. Let $A \subseteq Ct_{\mathbb{N}}(k)$ be a $\tilde{\Pi}^0_1$ -set.

- a) If $f : A \to \mathbb{R}$ is continuous, then f may be extended to a continuous $g : Ct_{\mathbb{N}}(k) \to \mathbb{R}$ such that the map $f \mapsto g$ is continuous.
- b) If $f : A \to \mathbb{N}$ is continuous, then f may be extended to a continuous $g : Ct_{\mathbb{N}}(k) \to \mathbb{N}$ such that the map $f \mapsto g$ is continuous.

Proof. Both a) and b) are trivial when k = 0 so assume that k > 0.

We prove a). The proof of b) is similar, but simpler.

Let $\bar{\pi}_k : Ct_{\mathbb{N}}(k) \to Ct_{\mathbb{R}}^E(k)$ be as in Corollary 3.2. Let *B* be the image of *A* under $\bar{\pi}_k$. If *x* is in the range of $\bar{\pi}_k$ we see from Corollary 3.2 that

$$x = \bar{\pi}_k(x \circ \bar{\pi}_{k-1}).$$

Then f' defined by $f'(x) = f(x \circ \overline{\pi}_{k-1})$ is a continuous map from B to \mathbb{R} , continuously depending on f.

By Proposition 2.16 a) there is a sequence $\{f'_n\}_{n\in\mathbb{N}}$ from $Ct^E_{\mathbb{R}}(k+1)$ such that whenever $x \in B$, $\{x_n\}_{n\in\mathbb{N}}$ is a sequence from $Ct^E_{\mathbb{R}}(k)$ and $x = \lim_{n\to\infty} x_n$, then $f'(x) = \lim_{n\to\infty} f'_n(x_n)$.

Let $f_n(z) = f'_n(\bar{\pi}_k(z)).$

Let $A = \bigcap_{n \in \mathbb{N}} A_n$ where each A_n is clopen.

Let

- g(z) = f(z) if $z \in A$.
- $g(z) = f_n(z)$ for the least n such that $z \notin A_n$ otherwise.

Then g is a continuous extension of f.

The construction of g from f is by composing continuous operators, so g depends continuously on f.

In order to prove b) we use part b) of Proposition 2.16 in a similar way.

This ends the proof of Theorem 3.6.

Corollary 3.7. If $A \subseteq Ct_{\mathbb{N}}(k)$ is $\tilde{\Pi}_{1}^{0}$, then $A \to \mathbb{N}$ is homeomorphic to a $\tilde{\Pi}_{1}^{0}$ -subset of $Ct_{\mathbb{N}}(k+1)$.

Proof. Let $f : A \to \mathbb{N}$. By Theorem 3.6 there will be an extension $f_1 : Ct(k) \to \mathbb{N}$ of f, continuous in f. Clearly, if $f_1 \in Ct_{\mathbb{N}}(k+1)$, then the restriction of f_1 to A is continuous in f_1 , so $A \to \mathbb{N}$ and

$$B = \{ f_1 \in Ct_{\mathbb{N}}(k+1) ; f \in A \to \mathbb{N} \}$$

are homeomorphic.

Let $\{x_i\}_{i\in\mathbb{N}}$ be a dense subset of $Ct_{\mathbb{N}}(k)$. Then

$$g \in B \Leftrightarrow \forall i(g(x_i) = (g \upharpoonright A)_1(x_i))$$

and this is Π_1^0 .

3.2. Some machinery.

Definition 3.8. For each $k \ge 0$ and $a \in Ct_{\mathbb{N}}(k)$ we define the *n*'th approximation a_n to *a* as follows:

For k = 0 we let $a_n = a$ if $a \le n$ and $a_n = 0$ if n < a. For k > 0 we let $a_n(x) = (a(x_n))_n$.

Lemma 3.9. (Essentially Grilliot [9])

For each $k \in \mathbb{N}$ and $a \in Ct_{\mathbb{N}}(k)$, we have that $a = \lim_{n \to \infty} a_n$.

For the sake of completeness, we give the proof. The point is that along with the proof, we give an algorithm for a *modulus of convergence* uniformly in the given a, i.e. when $k \ge 1$ we will give the algorithm for a map

$$M_k: Ct_{\mathbb{N}}(k-1) \times Ct_{\mathbb{N}} \to \mathbb{N}$$

such that

$$\forall a \in Ct_{\mathbb{N}}(k) \forall b \in Ct_{\mathbb{N}}(k-1) \forall m \in \mathbb{N}(m \ge M_k(b,a) \to a_m(b) = a(b)).$$

By Proposition 2.9 the existence of this modulus suffices to prove the lemma. We also define the modulus $M_0 : \mathbb{N} \to \mathbb{N}$.

In this proof we will observe the following conventions:

n, i, j etc. will denote natural numbers. f, g etc. will denote functions, or functionals one type below the type in question. F, G etc. will denote functionals of the type k in question, when $k \ge 2$. We will use induction on k.

Proof. k = 0: Clearly $\lim_{n \to \infty} i_n = i$ with modulus i, i.e. $n \ge i \Rightarrow i_n = i$. Thus we let $M_0(i) = i$

k = 1: Clearly $\lim_{n \to \infty} f_n(i) = \lim_{n \to \infty} (f(i_n))_n = f(i)$ with modulus $g(i) = \max\{i, f(i)\}$. Thus we let $M_1(i, f) = \max\{i, f(i)\}$.

k > 1: It is sufficient to show that $F(f) = \lim_{n \to \infty} F(f_n)$ and to compute a modulus G for this.

Convergency follows from the fact that F is continuous and that $f = \lim_{n \to \infty} f_n$.

Let g be the modulus for $f = \lim_{n \to \infty} f_n$ obtained by the induction hypothesis, i.e. $g(\xi) = M_{k-1}(\xi, f)$.

For each n and $\xi \in Ct_{\mathbb{N}}(k-2)$, let

- (1) $h_n(\xi) = f_m(\xi)$ for the least m such that $n \le m < g(\xi)$ and $F(f_m) \ne F(f)$ if there is such m.
- (2) $h_n(\xi) = f_{g(\xi)}(\xi)$ if there is no such m.

Then $h_n = f$ if there is no $m \ge n$ such that $F(f_m) \ne F(f)$, while $h_n = f_m$ for the least $m \ge n$ with $F(f_m) \ne F(f)$ otherwise.

Let $G(f) = max\{F(f), \mu n(F(h_n) = F(f))\}$. Then G will be the modulus for $F = \lim_{n \to \infty} F_n$.

Clearly h_n and G are computable as functions of n, f and F, so we let $M_k(f, F) = G(f)$ as defined above.

Lemma 3.10. If $a \in Ct_{\mathbb{N}}(k)$ and $n, m \in \mathbb{N}$, then $((a)_n)_m = a_{\min\{n,m\}}$.

This is proved by a trivial induction on k. Let $X_n^k = \{a_n ; a \in Ct_{\mathbb{N}}(k)\}.$

Lemma 3.11. Each X_n^k is a finite set.

Proof. We use induction on k:

 $X_n^0 = \{0, \dots, n\}.$

If $x_1 = (f_1)_n$, $x_2 = (f_2)_n$ and $x_1(y) = x_2(y)$ for all $y \in X_n^{k-1}$, then, using Lemma 3.10, we have for all $\xi \in Ct_{\mathbb{N}}(k-1)$

$$x_1(\xi) = (x_1)_n(\xi) = (x_1(\xi_n))_n = (x_2(\xi_n))_n = (x_2)_n(\xi) = x_2(\xi),$$

so $x_1 = x_2$. Since $(f)_n$ is bounded by n, we have an embedding of X_n^k into the finite set $X_n^{k-1} \to \{0, \ldots, n\}$. This embedding is actually onto.

The definition of the *n*'th approximation makes perfect sense for $a \in N(k)$ as well, with $(\perp)_n = \perp$. We then have

Lemma 3.12. Let $a \in N(k)$ be compact.

Then there is an $n_a \in \mathbb{N}$ such that for $n \geq n_a$ we have that $a \sqsubseteq (a)_n$.

Proof. We use induction on k. For k = 0 this is trivial, so let a be a compact element of N(k+1).

Then there are compact elements b_1, \ldots, b_r in N(k) and numbers m_1, \ldots, m_r such that a is minimal with the property that $a(b_i) = m_i$ for $i = 1, \ldots, r$.

Let $n_a = max\{n_{b_i}, m_i ; i \le r\}.$

 n_a will be the maximal value found in a or any of the hereditary sub-elements of a.

3.3. The construction. We will now construct the maps π_k by recursion on k. For k = 0 and k = 1 we will give explicit definitions. For k > 1 we will assume that π_{k-2} is defined and satisfies the requirements of the theorem.

For k > 1 the definition of π_k will for the sake of convenience be restricted to N(k). The definition is split into two cases, Case 1 and Case 2. In Case 2, the construction is easily extended to N(k), just interpreting the algorithm given over the partial objects as well. The construction in Case 2 will be effective. An important part of the proof will be to show that the π_k will be continuous on $\bar{N}(k)$. What we really do in this argument is to extend the part of π_k that is defined under Case 1 in an effective way to a partial continuous object consistent with the part of π_k constructed under Case 2. In proving consistency we rely on the fact that the total objects are dense for each space under consideration. The join of these two (the constructed part under Case 2 and the extension of the part constructed under Case 1) will finally form our π_k .

Discussing the consequences of the theorem, we pointed out that each π_k will have a partial inverse π_k^{-1} . Our separation in the two cases is needed in order to handle the problem that there is no total such inverse, due to the different topological nature of the Kleene-Kreisel functionals and the $Ct_{\mathbb{R}}^{E}$ -hierarchy, where each space is path connected. In Case 1 we will use the partial inverse that will exist in this case, and in Case 2 we want in a continuous way to bridge the gaps in the construction under Case 1.

In order to avoid too much notation, we will occasionally view X_n^k as a subset of $\bar{N}(k)$ instead of $Ct_{\mathbb{N}}(k)$.

Let $\pi_0(n) = n$ seen as an element of E(0).

Instead of proving the induction step, we prove a slightly stronger statement that we will need in Section 5.

Lemma 3.13. Let $k \geq 1$ and assume that π_0, \ldots, π_{k-1} are constructed according to the specifications of Theorem 3.1.

Then there is a continuous and total

$$\Pi_k : (N(k-1) \to E(0)) \to E(k)$$

such that $f = \prod_k (f) \circ \pi_{k-1}$ whenever $f : N(k-1) \to E(0)$ is total.

We will obtain $\pi_k(f)$ for $f \in N(k)$ by first modifying f to a $g: N(k-1) \to E(0)$ via the inclusion $N(0) \to E(0)$, and then use Π_k . It is of course sufficient to show that Π_k is continuous.

Proof of Lemma 3.13. There will be one direct construction for k = 1, and one depending on π_{k-2} for k > 1.

Let $\Pi_1 : (\mathbb{N}_\perp \to E(0)) \times E(0) \to E(0)$ be continuous such that for total $f \in \mathbb{N}_\perp \to E(0)$ and total $x \in E(0)$ we have

- $\Pi_1(f)(x) = f(0)$ if $x \le 0$.
- $\Pi_1(f)(x) = (1-y)f(n) + yf(n+1)$ when x = n+y and $0 \le y \le 1$.

We may choose Π_1 to be definable in *RealPCF*. From now on, let $k \geq 2$, let F be a total map in $N(k-1) \to E(0)$ and let $g \in \overline{E}(k-1)$. We will define $\Pi_k(F)(g) \in E(0)$ and prove that Π_k is continuous.

We will use ξ and η for elements in $\overline{N}(k-2)$. Let $\{\eta_n\}_{n\in\mathbb{N}}$ be an effectively enumerated dense subset of N(k-2).

We will separate the definition of $\Pi_k(F)(g)$ into two cases, and prove continuity later.

Let \mathbb{N}_E be the set of elements in E(0) representing natural numbers and let *nat* : $E(0) \rightarrow N(0)$ send representatives of n to n.

Case 1: $g(\pi_{k-2}(\eta_n)) \in \mathbb{N}_E$ for all n.

By continuity and the totality of g we have that $g(\pi_{k-2}(\xi)) \in \mathbb{N}_E$ for all $\xi \in \overline{N}(k-2)$. Let

$$f_g = \lambda \xi \in N(k-2).nat(g(\pi_{k-2}(\xi))).$$

Then $f_g \in \overline{N}(k-1)$ and we let $\Pi_k(F)(g) = \pi_0(F(f_g))$. It is at this point that we ensure that $\pi_k(F)(\pi_{k-1}(f)) = \pi_0(F(f))$, see the end of the proof.

Case 2: Otherwise.

In order to save notation and making the construction more transparent, we behave as if we operate over \mathbb{R} and with $\mathbb{N} \subset \mathbb{R}$. As mentioned above, if we view this definition as an algorithm for exact computations over the partial reals, we actually define an effective map $\Pi'_k \in (N(k-1) \to E(0)) \to (E(k-1) \to E(0))$ in this case, and we will let $\Pi'_k \sqsubseteq \Pi_k$ in the end.

Let d(g,n) be the distance from $g(\pi_{k-2}(\eta_n))$ to \mathbb{N} . Since g is continuous, there is an $\epsilon > 0$ and infinitely many n such that $d(g, n) > \epsilon$. Thus

$$\sum_{n=0}^{\infty} d(g,n) = \infty$$

Let

•
$$z_n(g) = 1$$
 if $\sum_{i < n} d(g, i) \le 1$

•
$$z_n(g) = 0$$
 if $\sum_{i < n} d(g, i) > 1$.

•
$$z_n(g) = y$$
 such that $\sum_{i < n} d(g, i) + y = 1$ otherwise

Since g is total, this makes sense.

Each $x \in \mathbb{R}$ will induce a probability distribution μ_x on \mathbb{N} by

- If $x \le 0$, then $\mu_x(0) = 1$ and $\mu_x(n) = 0$ for n > 0.
- If $n \frac{1}{3} \le x \le n + \frac{1}{3}$, then $\mu_x(n) = 1$ and $\mu_x(m) = 0$ for $m \ne n$. If $n + \frac{1}{3} \le x \le n + \frac{2}{3}$, let $y \in [0, 1]$ be such that $x = n + \frac{1+y}{3}$. Then let $\mu_x(n) = 1 y$, $\mu_x(n+1) = y$ and $\mu_x(m) = 0$ for all other m.
- μ_x will induce a probability distribution on X_n^{k-1} as follows:

For $a \in X_n^{k-1}$, let

$$\mu_{n,g}(a) = \prod_{b \in X_n^{k-2}} \mu_{\min\{n,g(\pi_{k-2}(b))\}}(a(b)).$$

Claim 1:

$$\sum_{a \in X_n^{k-1}} \mu_{n,g}(a) = 1$$

Proof. Each function $h: X_n^{k-2} \to \{0, \ldots, n\}$ corresponds to one and only one $a \in X_n^{k-1}$. $\mu_{\min\{n,g(\pi_{k-2}(b))\}}$ is a probability distribution on $\{0, \ldots, n\}$, so $\mu_{n,g}$ can be viewed as

the product distribution. Claim 1 follows.

Now, let

$$\Pi_k(F)(g) = \sum_{n \in \mathbb{N}} \left(d(g, n) \cdot z_n(g) \cdot \sum_{a \in X_n^{k-1}} (F(a) \cdot \mu_{n,g}(a)) \right).$$

This ends the construction in Case 2.

It is easy to see that the constructions in Cases 1 and 2 are continuous separately. Moreover, the domain for Case 2 is open. In order to prove the continuity of $\Pi_k(F)$ and of the map Π_k it is sufficient to show that if $g \in \overline{R}(k-1)$ falls under Case 1 and $\epsilon > 0$ is given, there are compact approximations δ and τ to F and g resp. such that for any total $F' \in N(k-1) \to E(0)$ extending δ and any total $g' \in E(k-1)$ extending τ we have that

$$|\Pi_k(F')(g') - F(g)| < \epsilon.$$

So, let F, g and $\epsilon > 0$ be given as above, and without loss of generality, assume that $\epsilon < 1$. Let $f \in N(k-1)$ be defined by

$$f(\xi) = nat(g(\pi_{k-2}(\xi)))$$

Claim 2: There are $\sigma \sqsubseteq f$, $\delta \sqsubseteq F$ and $n_0 \in \mathbb{N}$ such that $\delta(\sigma)$ has length $< \frac{\epsilon}{3}$, such that for all $n \ge n_0$ and f' extending σ we have that $\sigma \sqsubseteq f'_n$, such that $\sigma(\tau) < n_0$ whenever defined and such that if $n \leq n_0$ and $a \in X_n^{k-1}$ then $\delta(a)$ has length $\leq \frac{\epsilon}{3}$.

Proof. First pick $\sigma \sqsubseteq f$ and $\delta' \sqsubseteq F$ such that $\delta'(\sigma)$ has length $< \frac{\epsilon}{3}$.

By Lemma 3.12 there is an n_0 such that $(\sigma)_n = \sigma$ for $n \ge n_0$ and such that n_0 exceeds all values of σ .

Given n_0 , we may find $\delta \sqsubseteq F$ with $\delta' \sqsubseteq \delta$ such that $\delta(a)$ has length $\leq \frac{\epsilon}{3}$ for all $a \in X_n^{k-1}$ with $n \leq n_0$.

This ends the proof of the claim.

Let $\hat{\sigma} = \sigma \pm \frac{1}{3}$, i.e. σ is the compact in $N(k-2) \to E(0)$ where each value $c \in \mathbb{N}$ is replaced by the interval $[c - \frac{1}{3}, c + \frac{1}{3}]$. Thus $\hat{\sigma} \sqsubseteq \lambda \xi \in N(k-2).g(\pi_{k-2}(\xi))$. Let $\tau_0 \sqsubseteq g$ be compact such that

$$\hat{\sigma} \sqsubseteq \lambda \xi \in N(k-2).\tau_0(\pi_{k-2}(\xi)).$$

Let

$$M = max\{|x| ; x \in \delta(a) \land a \in X_n^{k-1} \land n \le n_0\}.$$

Let $\tau_1 \sqsubseteq g$ be such that for any total g' extending τ_1 we have that

$$\sum_{n \le n_0} d(g', n) < \frac{\epsilon}{3M}$$

We may let $\tau_0 \sqsubseteq \tau_1$.

We complete the proof of the continuity by showing Claim 3: If F' and g' are total extensions of δ and τ_1 resp., then

 $|\Pi_k(F')(g') - F(f)| < \epsilon.$

Proof. The proof will be divided into the same cases as the construction.

Case 1: $g'(\pi_{k-2}(\xi)) \in \mathbb{N}_E$ for all $\xi \in \overline{N}(k-2)$.

Then let $f' = \lambda \xi \in N(k-2).nat(g'(\pi_{k-2}(\xi))).$

By the choice of $\tau_0 \sqsubseteq \tau_1$ we have that $\sigma \sqsubseteq f'$. Since $\delta(\sigma)$ has length $< \frac{\epsilon}{3}$ and since $\Pi_k(F')(g') = F'(f')$ we actually have that $|\Pi_k(F')(g') - F(f)| < \frac{\epsilon}{3}$.

Case 2: Otherwise.

Let n_1 be minimal such that

$$\sum_{n \le n_1} d(g', n) \ge 1.$$

Since $\sum_{n \leq n_0} d(g', n) < \frac{\epsilon}{3M}$, $\epsilon < 1$ and $M \geq 1$, we see that $n_0 < n_1$. Then

$$\Pi_{k}(F')(g') = \sum_{n \le n_{0}} \left(d(g', n) \cdot \sum_{a \in X_{n}^{k-1}} (F'(a) \cdot \mu_{n,g'}(a)) \right) + \sum_{n_{0} < n \le n_{1}} \left(d(g', n) \cdot z_{n}(g') \cdot \sum_{a \in X_{n}^{k-1}} (F'(a) \cdot \mu_{n,g'}(a)) \right).$$

Since $\sum_{n \leq n_0} d(g', n) < \frac{\epsilon}{3M}$ and $|F'(a)| \leq M$ whenever $n \leq n_0$ and $a \in X_n^{k-1}$, the first part will be bounded by $\frac{\epsilon}{3}$.

Subclaim 3.1: If $n > n_0$, then

$$\mu_{n,g'}(\{a \in X_n^{k-1}; |F'(a) - F(f)| < \frac{\epsilon}{3}\}) = 1.$$

Proof of subclaim: Let $n > n_0$ and $a \in X_n^{k-1}$. We will show that we either have that $\mu_{n,g'}(a) = 0$ or that $|F'(a) - F(f)| < \frac{\epsilon}{3}$.

 $\begin{aligned} \mu_{n,g'}(a) &= 0 \text{ or that } |F'(a) - F(f)| < \frac{\epsilon}{3}. \\ Subcase \ 1. \text{ For some } b \in X_n^{k-2} \text{ we have that } \sigma(b) \text{ is defined and } \sigma(b) \neq a(b). \\ \text{Then } \sigma(b) < n_0 \text{ by choice of } n_0. \text{ Since } g' \text{ extends } \tau_0, \text{ we have that } \end{aligned}$

$$u(0) < n_0$$
 by choice of n_0 . Since g extends n_0 , we have that

$$g'(\pi_{k-1}(b)) \in [\sigma(b) - \frac{1}{3}, \sigma(b) + \frac{1}{3}],$$

so the distance from a(b) to $min\{n_0, g'(\pi_{k-2}(b))\}$ is at least $\frac{2}{3}$. It follows that $\mu_{n,g'}(a) = 0$. Subcase 2. Otherwise.

Then $\sigma(b) = a(b)$ whenever $\sigma(b) \in \mathbb{N}$ and we may extend σ to σ_1 such that $\sigma_1(b) = a(b)$ for all $b \in X_n^{k-2}$ and $\sigma_1(\xi) \leq n$ for all ξ . Let f' be a total extension of σ_1 such that $f'(\xi) \leq n$ for all total ξ . Then by Lemma 3.10

$$f'_n(\xi) = f'(\xi_n) = \sigma_1(\xi_n) = a(\xi_n) = a_n(\xi) = a(\xi)$$

so $f'_n = a$.

Thus $a = f'_n$ for some f' extending σ , and by the choice of σ and n_0 , we have that $\sigma \sqsubseteq a$. Since $\sigma \sqsubseteq f$, $\delta(\sigma)$ has length $< \frac{\epsilon}{3}$ and $\delta \sqsubseteq F \sqcap F'$ it follows that

$$|F'(a) - F(f)| < \frac{\epsilon}{3}$$

This ends the proof of the subclaim. Continuity of the construction will then follow from *Subclaim 3.2:*

$$\sum_{n_0 < n \le n_1} \left(d(g', n) \cdot z_n(g') \cdot \sum_{a \in X_n^{k-1}} (F'(a) \cdot \mu_{n,g'}(a)) \right) - F(f) \right| < \frac{2\epsilon}{3}.$$

Proof of subclaim: Since $\tau_1 \sqsubseteq g'$, we have that

$$\sum_{n_0 < n \le n_1} d(g', n) \cdot z_n(g') > 1 - \frac{\epsilon}{3M}.$$

It follows that

$$\left| \sum_{n_0 < n \le n_1} \left(d(g', n) \cdot z_n(g') \cdot \sum_{a \in X_n^{k-1}} (F'(a) \cdot \mu_{n,g'}(a)) \right) - F(f) \right| \le \frac{\epsilon}{3M} \cdot |F(f)| + \sum_{n_0 < n \le n_1} \left(d(g', n) \cdot z_n(g') \cdot \sum_{a \in X_n^{k-1}} (|F'(a) - F(f)| \cdot \mu_{n,g'}(a)) \right)$$

We have that $|F(f)| \leq M$, and by subclaim 3.1 the above is bounded by

$$\frac{\epsilon}{3} + \sum_{n_0 < n \le n_1} d(g', n) \cdot z_n(g') \cdot \frac{\epsilon}{3} \le \frac{\epsilon}{3} + \frac{\epsilon}{3} = \frac{2\epsilon}{3}.$$

This ends the proof of subclaim 3.2, of claim 3 and of the continuity of the construction. Thus Lemma 3.13 is proved.

We may now end the proof of Theorem 3.1, where Lemma 3.13 provides us with the induction step. Let $F_0 \in \overline{N}(k)$. Let F be the corresponding total $F \in N(k-1) \to E(0)$ and let $\pi_k(F_0) = \prod_k(F)$.

We show that $\pi_k(F)(\pi_{k-1}(f)) = \pi_0(F(f))$ by induction on k. For k = 1 this is trivial. For k > 1, let $f \in \overline{N}(k-1)$ and let $g = \pi_{k-1}(f)$.

By the induction hypothesis, the f_g constructed from g in Case 1 will be the f given. Then $\Pi_k(F)(g) = \pi_0(F(f))$, i.e.

$$\pi_k(F)(\pi_{k-1}(f)) = \pi_0(F(f)).$$

In the proof of the continuity we started with a total F, a total g and some $\epsilon > 0$ and showed the existence of approximations demonstrating the continuity of Π_k . Now the set of triples (δ, τ, ϵ) that are constructed in this proof will be decidable, and the part of Π_k that can be constructed from this set of triples will be effective. Thus Π_k will be the join of two effective partial functionals, and thus it will itself be effective.

Remark 3.14. Note that the construction under Case 2 will not terminate for g that falls under Case 1. We use this to observe that if $F \in N(k)$ is not total, we may let $f \in \overline{N}(k-1)$ be such that $F(f) = \bot$. Then $\pi_k(F)(\pi_{k-1}(f))$ will be undefined, because neither Case 1 nor Case 2 will provide us with a value. We use this to prove Corollary 3.5

4. The intensional functionals

An alternative approach to higher type objects over \mathbb{R} is based on representations of the reals using intensional objects. A similar hierarchy was studied in Bauer, Escardó and Simpson [3]. We gave a full treatment in Normann [17]. We call this hierarchy the *I*-hierarchy, where *I* stands for 'intensional'.

Definition 4.1. Let $ab_1b_2\cdots$ be an element in $\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$, which we view as a set of functions defined on \mathbb{N} .

Let

$$\rho_0^I(ab_1b_2\cdots) = a + \sum_{n>0} b_n \cdot 2^{-n}$$

Let $Ct^{I}_{\mathbb{R}}(0) = \mathbb{R}$.

Definition 4.2. Let I(0) be the algebraic domain consisting of the empty sequence e, all finite sequences $ab_1 \dots b_n$ and all infinite sequences $ab_1b_2 \dots$, where in the two latter cases $a \in \mathbb{Z}$ and each $b_i \in \{-1, 0, 1\}$.

I(0) is ordered by sequence end-extensions.

As we have seen, each maximal element in I(0) will determine a real via ρ_0^I . We let $\overline{I}(0)$ be the set of maximal elements in I(0), and we let

$$ab_1b_2\cdots\approx_0^I cd_1d_2\cdots$$

if they represent the same real.

We now extend these concepts to higher types in analogy with the constructions for the N-hierarchy and the E-hierarchy:

Definition 4.3. By recursion on k we let

- a) $I(k+1) = I(k) \rightarrow I(0)$ in the category of algebraic domains.
- b) If x_1 and x_2 are in I(k+1), we let $x_1 \approx_{k+1}^{I} x_2$ if for all $y_1, y_2 \in I(k)$, if $y_1 \approx_{k}^{I} y_2$, then $x_1(y_1) \approx_{0}^{I} x_2(y_2)$.

- c) Let $\overline{I}(k+1) = \{x \in I(k+1) ; x \approx_{k+1} x\}.$ d) Let ρ_{k+1}^I map an element x of $\overline{I}(k+1)$ to a function $\rho_{k+1}^I(x) : Ct_{\mathbb{R}}^I(k) \to \mathbb{R}$ defined as follows: $(r)(o_i^I(u)) = o_0^I(u)$

$$\begin{split} \rho^{I}_{k+1}(x)(\rho^{I}_{k}(y)) &= \rho^{I}_{0}(x(y)).\\ \text{Let } Ct^{I}_{\mathbb{R}}(k+1) &= \{\rho^{I}_{k+1}(x) \ ; \ x \in \bar{I}(k+1)\}. \end{split}$$

One motivation for using the *I*-hierarchy is that whenever $A \subseteq \mathbb{N}^{\mathbb{N}}$ and $f : A \to \mathbb{R}$ is continuous, then there is an $\hat{f} : A \to \bar{I}(0)$ such that $f(x) = \rho_0^I(\hat{f}(x))$ for each $x \in A$.

Definition 4.4. Let T be a topological space.

Let $\mathcal{R}(T)$ be the subtopology where the open sets are

 $\{f^{-1}(O) ; f: T \to \mathbb{R} \text{ is continuous and } O \subseteq \mathbb{R} \text{ is open } \}.$

Clearly every clopen set in T will be clopen in $\mathcal{R}(T)$, i.e. $\mathcal{Z}(T) \subseteq \mathcal{R}(T)$. The following is essentially observed by Bauer, Escardó and Simpson [3]:

Proposition 4.5. Let T be a topological space that is hereditarily Lindelöf. Then the following are equivalent:

- a) For every continuous $f: T \to \mathbb{R}$ there is a continuous $\hat{f}: T \to \overline{I}(0)$ such that $f(x) = \rho_0^I(\hat{f}(x))$ for all $x \in T$.
- b) $\mathcal{R}(T)$ is zero-dimensional (i.e. has a basis of clopen sets).

The key problem under discussion is:

Is
$$Ct^{E}_{\mathbb{R}}(k) = Ct^{I}_{\mathbb{R}}(k)$$
 for a given k?

For $k \leq 2$, the equality was proved in [3]. They also proved that the statement that

$$Ct^E_{\mathbb{R}}(3) = Ct^I_{\mathbb{R}}(3)$$

followed from the assumption that $Ct_{\mathbb{N}}(2)$ is zero-dimensional.

We will extend these results. As a tool, we will use the S-hierarchy introduced in Normann [17]. The S-hierarchy is not as natural as the E-hierarchy and the I-hierarchy, but it too is equipped with hereditarily total elements $\overline{S}(k)$, and the extensional collapses $Ct^{S}_{\mathbb{D}}(k)$. The S-hierarchy is in some sense a smoothened I-hierarchy. We will give the technical definitions below.

Definition 4.6. By recursion on k we define the domain S(k) and the binary relation \sim_k on S(k) as follows

- S(0) = I(0) with the same ordering. $x_1 \sim_0 x_2$ if there are maximal extensions y_1 of x_1 and y_2 of x_2 such that $\rho_0^I(y_1) =$ $\rho_0^I(y_2).$
- S(k+1) is the set of Scott-continuous functions $f: S(k) \to S(0)$ such that

$$x_1 \sim_k x_2 \Rightarrow f(x_1) \sim_0 f(x_2).$$

S(k+1) is ordered by the pointwise ordering.

If $f_1 \in S(k+1)$ and $f_2 \in S(k+1)$, we let $f_1 \sim_{k+1} f_2$ if

$$\forall x_1 \in S(k) \forall x_2 \in S(k) (x_1 \sim_k x_2 \Rightarrow f_1(x_1) \sim_0 f_2(x_2)).$$

By construction, S(0) is an algebraic domain. In Normann [17] it is proved that S(k+1)is a closed subset of $S(k) \to S(0)$ and, with the restricted ordering, is an algebraic domain. The compact objects in the sense of S(k+1) are the compact objects in the sense of $S(k) \to S(0)$ that are in S(k+1), but the boundedness relation is in general not the same.

The relation \sim_k is reflexive and symmetric, but not transitive.

Definition 4.7. We let $\overline{S}(0)$ be the maximal elements with the partial equivalence relation

 \approx_0^S which will be \sim_0 restricted to the maximal objects. We let $\rho_0^S = \rho_0^I$. By recursion on k we then define a partial equivalence relation \approx_k^S on S(k) for each k in analogy with our previous constructions of hierarchies. Let $\bar{S}(k) = \{x \in S(k) ; x \approx_k^S x\}$, and define ρ_k^S in analogy with ρ_k^I . We then define $Ct_{\mathbb{R}}^S(k)$ in analogy with $Ct_{\mathbb{R}}^I(k)$.

Proposition 4.8. (Normann [17])

- a) Uniformly in any compact element p in S(k) there is an extension to an element $\xi(p) \in S(k).$
- b) For each k, $Ct^S_{\mathbb{R}}(k) = Ct^I_{\mathbb{R}}(k)$.

Proposition 4.9. (Normann [17]) If $x_1, x_2 \in \overline{S}(k)$, then

$$x_1 \approx^S_k x_2 \Leftrightarrow x_1 \sim_k x_2.$$

Moreover, \sim_k is a closed relation on S(k).

Finally, if y_1 and y_2 are compacts in S(k) such that $y_1 \sim_k y_2$, then there are total extensions x_1 and x_2 of y_1 and y_2 resp. such that $x_1 \approx_k^S x_2$.

Our aim is to show that the hierarchies $\{Ct^{I}_{\mathbb{R}}(k)\}_{k\in\mathbb{N}}$ and $\{Ct^{E}_{\mathbb{R}}(k)\}_{k\in\mathbb{N}}$ are identical, assuming that $\mathcal{R}(Ct_{\mathbb{N}}(k))$ is zero-dimensional for each k. In order to make use of this assumption, we will consider the S-hierarchy and the quotient space of $\bar{S}(k)$ under the consistency relation. This will be an intermediate stage between $\bar{S}(k)$ and $\bar{S}(k)/\approx_k^S$, an intermediate stage that will enjoy some of the topological qualities of $Ct_{\mathbb{N}}(k)$. (See Lemma 4.14 below.)

Definition 4.10. Let $x_1, x_2 \in \overline{S}(k)$. We define the relation C_k on $\overline{S}(k)$ by: $C_k(x_1, x_2) \Leftrightarrow x_1$ and x_2 are consistent, i.e. bounded in S(k).

Lemma 4.11. C_k is an equivalence relation on $\overline{S}(k)$.

Proof. C_0 is the identity relation on $\overline{S}(0)$.

By Proposition 4.8 a), two higher type total objects in the S-hierarchy are consistent if and only if they are identical when restricted to total inputs. This defines an equivalence relation.

Definition 4.12. Let $T(k) = \overline{S}(k)/C_k$ with the quotient topology.

Recall that the topology on T(k) then will be sequential.

Lemma 4.13.

- a) T(0) is homeomorphic to $\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$.
- b) T(k+1) is homeomorphic to a closed subspace of $T(k) \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$.

Proof. a) is trivial. In order to prove b), we observe that $\overline{S}(k+1)$ will consist of the total elements in $S(k) \to S(0)$ that send extensionally equivalent elements in $\bar{S}(k)$ to extensionally equivalent elements of S(0).

Being extensionally equivalent is a closed relation on $\bar{S}(0)$, so $\bar{S}(k+1)$ will be a closed subset of the set of total elements from $\bar{S}(k)$ to $\bar{S}(0)$.

 C_{k+1} is just the restriction of the equivalence relation of consistency for total elements on $S(k) \to S(0)$, and $\bar{S}(k+1)$ will consist of full equivalence classes for this consistency relation.

We then obtain the lemma from Lemma 2.8.

Lemma 4.14. If $k \geq 1$, then T(k) is homeomorphic to a $\tilde{\Pi}_1^0$ -subspace of $Ct_{\mathbb{N}}(k)$.

Proof. For k = 1, this is proved in [3]. Our proof is inspired by the proof in [3], but we give a slightly different proof in order to prepare for the induction step.

If $f \in T(1)$, then f is (the equivalence class representing) a total map from $\mathbb{Z} \times \{-1,0,1\}^{\mathbb{N}^+}$ to $\mathbb{Z} \times \{-1,0,1\}^{\mathbb{N}^+}$.

 $\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$ is homeomorphic to a σ -compact subset of $\mathbb{N}^{\mathbb{N}}$, i.e. a countable union of compact sets.

For $a \in \mathbb{Z}$ and $g \in \{-1, 0, 1\}^{\mathbb{N}^+}$ let $f_{a,i}(g)$ be the *i*'th element in the sequence f(a * g) (where * is concatenation between a finite sequence and a function). Using the fan functional, we may find a number k(f, a, i) that codes the behavior of $f_{a,i}$ on $\{-1, 0, 1\}^{\mathbb{N}^+}$ and such that $f_{a,i}$ is recoverable from k(f, a, i).

If $\{a_n\}_{n\in\mathbb{N}}$ is a 1-1 enumeration of \mathbb{Z} , we let $c_f(n,i) = k(f,a_n,i)$, and we let

$$X_1 = \{ c_f \; ; \; f \in T(1) \},\$$

then X_1 is homeomorphic to T(1).

We call c_f the code for f.

In order to see that $X_1 \in \tilde{\Pi}_1^0$, we observe that the $\tilde{\Pi}_1^0$ -subsets of $Ct_{\mathbb{N}}(1)$ are exactly the closed subsets, and further that

(1) The set of codes in $\mathbb{N}^{\mathbb{N}}$ for total elements in

$$\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+} \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$$

is closed.

(2) The set of codes for total elements in

$$\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+} \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$$

that represents elements in T(1) is closed.

(1) is trivial. (2) is seen as follows: Let $\{(\sigma_n, \tau_n)\}_{n \in \mathbb{N}}$ be an enumeration of all pairs (σ, τ) of compact elements in S(0) such that $\sigma \sim_0 \tau$. By the third part of Proposition 4.9, let ξ_n and η_n be equivalent, total extensions of σ_n and τ_n . Then

$$f: \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+} \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$$

represents an element in T(1) if and only if

$$\forall n(f(\xi_n) \sim_0 f(\eta_n))$$

if and only if

$$\forall n \forall i (\overline{f(\xi_n)}(i) \sim_0 \overline{f(\eta_n)}(i),$$

where $\overline{g}(i) = (g(0), \dots, g(i-1))$ whenever g is defined on N.

The matrix defines a clopen set, so we are through with the induction start.

Now assume that T(k) is homeomorphic to a Π_1^0 -subset X_k of $Ct_{\mathbb{N}}(k)$.

By Lemma 4.13, T(k + 1) is homeomorphic to a closed subset A of $T(k) \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$, and using the argument for 2. under the case for k = 1, we see that A is indeed a countable intersection of clopen sets.

Clearly $T(k) \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$ is homeomorphic to $X_k \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$. Furthermore, $\mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$ is homeomorphic to a $\tilde{\Pi}_1^0$ -subset of $\mathbb{N}^{\mathbb{N}}$, and $X_k \to \mathbb{N}^{\mathbb{N}}$ is homeomorphic to $X_k \times \mathbb{N} \to \mathbb{N}$.

By Corollary 3.7, we see that $X_k \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$ then is homeomorphic to a $\tilde{\Pi}_1^0$ -subset of $Ct_{\mathbb{N}}(k+1)$.

If we use all these homeomorphisms to map $A \subseteq T(k) \to \mathbb{Z} \times \{-1, 0, 1\}^{\mathbb{N}^+}$ to $B \subseteq Ct_{\mathbb{N}}(k+1)$, we see that B will be $\tilde{\Pi}_1^0$ and homeomorphic to T(k+1). This ends the proof of the lemma.

Remark 4.15. With some care, we may prove that the sets are Π_1^0 , i.e. they will be the intersection of an effectively given sequence of clopen sets.

Lemma 4.16. Let (Ω, Ω) be the lazy natural numbers as defined in the proof of Lemma 2.13.

Then $\overline{\Omega} \times T(k)$ with the quotient topology is homeomorphic to a $\widetilde{\Pi}_1^0$ -subspace of $Ct_{\mathbb{N}}(k)$.

Proof. Clearly $\overline{\Omega} \times Ct_{\mathbb{N}}(k)$ is homeomorphic to a Π_1^0 -subset of $Ct_{\mathbb{N}}(k)$. Then the lemma follows from Lemma 4.14.

Now we are ready to prove

Theorem 4.17. If $Ct_{\mathbb{N}}(n)$ is zero-dimensional, then

$$Ct_{\mathbb{R}}^{S}(n+1) = Ct_{\mathbb{R}}^{E}(n+1)$$

as topological spaces.

Proof. By the assumption, $Ct_{\mathbb{N}}(k)$ is zero-dimensional for $k \leq n$, i.e. the assumption of the theorem holds for all $k \leq n$.

By a simultaneous induction on n satisfying the assumption, we will prove the following three claims:

- (1) $Ct^S_{\mathbb{R}}(n+1) = Ct^E_{\mathbb{R}}(n+1)$ as sets.
- (2) The quotient topology on $Ct^{S}_{\mathbb{R}}(n+1)$ coincides with the quotient topology on $Ct^{E}_{\mathbb{R}}(n+1)$.
- (3) The quotient topology on $\overline{\Omega} \times Ct^S_{\mathbb{R}}(n+1)$ coincides with the quotient topology on $\overline{\Omega} \times Ct^E_{\mathbb{R}}(n+1)$.

For n = -1 the three claims will hold, so we are on safe ground in proving this by inducton. Let $n \ge 1$ and assume that the three claims hold for m = n - 1.

Proof of (1). As remarked in Section 2, $Ct^E_{\mathbb{R}}(n+1)$ consists of exactly all continuous functions $F: Ct^E_{\mathbb{R}}(n) \to \mathbb{R}$. It follows from the induction hypothesis, (1) and (2), that $Ct^S_{\mathbb{R}}(n+1) \subseteq Ct^E_{\mathbb{R}}(n+1)$.

By Lemma 4.14 and the assumption it follows that T(n) is zero-dimensional, and in particular, $\mathcal{R}(T(n))$ is zero-dimensional. Thus, if $f \in Ct^E_{\mathbb{R}}(n+1)$, then $f : Ct^S_{\mathbb{R}}(n) \to \mathbb{R}$ is continuous, and it factors through a continuous $\hat{f} : T(n) \to \mathbb{R}$. It follows from Proposition 4.5 that $f \in Ct^S_{\mathbb{R}}$, and the equality of the two sets is established. *Proof of (2).* Since both $Ct^S_{\mathbb{R}}(n+1)$ and $Ct^E_{\mathbb{R}}(n+1)$ are sequential topological spaces, it is sufficient to show that the convergents sequences are the same for the two topologies.

This amounts to prove that $\overline{\Omega} \times Ct^E_{\mathbb{R}}(n) \to \mathbb{R}$ and $\overline{\Omega} \times Ct^S_{\mathbb{R}}(n) \to \mathbb{R}$ are identical as sets, where all topologies are the relevant quotient topologies.

Using Lemma 4.16 instead of Lemma 4.14, we may use the same argument as in (1).

Proof of (3). Following the line of thought from (2), we have to prove that $\overline{\Omega} \times Ct^E_{\mathbb{R}}(n) \to \mathbb{R}$ and $\overline{\Omega} \times Ct^S_{\mathbb{R}}(n) \to \mathbb{R}$ as topological spaces have the same convergent sequences. But in the quotient topologies of these products, the convergent sequences will be exactly the products of convergent sequences in the factors (using liftings of convergent sequences to see that the product of two convergent sequences is convergent), and then (3) follows from (2).

We also obtain the following:

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Theorem 4.18. Assume that $\mathcal{R}(Ct_{\mathbb{N}}(n))$ is zero-dimensional. Then

 $Ct^S_{\mathbb{R}}(n+1) = Ct^E_{\mathbb{R}}(n+1).$

Proof. First we observe that $\mathcal{R}(Ct_{\mathbb{N}}(k))$ is zero-dimensional for each $k \leq n$

Claim: If $k \leq n$ and $A \subseteq Ct_{\mathbb{N}}(k)$ is Π_1^0 , then $\mathcal{R}(A)$ is zero-dimensional.

Proof of claim: Let $f: A \to \mathbb{R}$ be continuous and $O \subseteq \mathbb{R}$ be open.

By Theorem 3.6 a), f may be extended to a continuous $g: Ct_{\mathbb{N}}(k) \to \mathbb{R}$. Then, by the assumption, $g^{-1}[O]$ will be the union of clopen sets in $Ct_{\mathbb{N}}(k)$, and thus $f^{-1}[O]$ will be the union of clopen sets in A.

This ends the proof of the claim.

From now on we may use the same argument as in the proof of Theorem 4.17.

Since $Ct^S_{\mathbb{R}}(k) = Ct^I_{\mathbb{R}}(k)$ with the same topology, we have proved that if $\mathcal{R}(Ct_{\mathbb{N}}(k))$ is zero-dimensional for each k, the extensional and intensional hierarchies will coincide.

5. A TOPOLOGICAL CHARACTERIZATION

We have only used Proposition 4.5 one way. We will now prove the converse of Theorem 4.18, using the other direction of Proposition 4.5.

As a tool, we will construct continuous maps $\pi_k^S : \overline{N}(k) \to \overline{S}(k)$ and embeddings $\overline{\pi}_k^S : Ct_{\mathbb{N}}(k) \to Ct_{\mathbb{R}}^S(k)$ using as far as possible the same construction as for π_k and $\overline{\pi}_k$ from the proofs of Theorem 3.1 and its corollaries.

In Normann [17] we proved that the hierarchies $\{Ct^E_{\mathbb{R}}(k)\}_{k\in\mathbb{N}}$ and $\{Ct^I_{\mathbb{R}}(k)\}_{k\in\mathbb{N}}$ have a maximal common core, i.e. there are maximal isomorphic sub-hierarchies of these type structures. Using the similarities of the constructions of $\bar{\pi}_k$ and $\bar{\pi}^s_k$ we obtain that the $Ct_{\mathbb{N}}$ -hierarchy can be embedded into this core, see Theorem 5.4.

Lemma 5.1. There is a total, continuous map, called the "normalizer",

 $norm: S(0) \rightarrow S(0)$

representing the identity map on \mathbb{R} such that whenever $x \in \overline{S}(0)$ and $\rho_0^S(x) = n \in \mathbb{N}$, then $norm(x) = n00\cdots$.

Proof. From finite information about x we may split between two overlapping cases:

(1) $\exists n(\rho_0^S(x) \in (n - \frac{1}{2}, n + \frac{1}{2}))$ (2) $\exists n(\rho_0^S(x) \in (n + \frac{1}{3}, n + \frac{2}{3})).$ In Case (2) we let norm(x) = x.

In Case (1) we find the k'th element in the sequence norm(x) by recursion on k as follows:

The first element is n (we are certain that n has an extension equivalent to x).

Assume that we at the k'th stage have decided that norm(x) starts with n and then k-1 zeros, and that we know that x represents a real in $(n-2^{-k}, n+2^{-k})$.

Then from a finite part of x we may separate between the overlapping cases

- x represents a real in $(n 2^{-(k+1)}, n + 2^{-(k+1)})$
- x represents a real in $(n-2^{-k}, n-2^{-(k+2)})$
- x represents a real in $(n + 2^{-(k+2)}, n + 2^{-k})$.

In the first case, we add a new zero to norm(x) and continue. In the other cases, we let norm(x) be an extension of $n * (0)^{k-1}$ equivalent to x, which we can find continuously in x.

Theorem 5.2. For each $k \in \mathbb{N}$ there is a continuous total map

$$\pi_k^S: N(k) \to S(k)$$

such that whenever $x \in \overline{N}(k+1)$ and $y \in \overline{N}(k)$ we have

$$\rho_0^N(x(y)) = \rho_0^S(\pi_{k+1}^S(x)(\pi_k^S(y)))$$

We will prove the theorem below. Let us first observe:

Corollary 5.3. For each $k \in \mathbb{N}$ there is an injective, continuous map

$$\bar{\pi}_k^S : Ct_{\mathbb{N}}(k) \to Ct_{\mathbb{R}}^S$$

such that

- i) $\bar{\pi}_0^S$ is the standard inclusion map from \mathbb{N} to \mathbb{R} . ii) For each $k \in \mathbb{N}$, each $F \in Ct_{\mathbb{N}}(k+1)$ and each $a \in Ct_{\mathbb{N}}(k)$ we have that

$$F(a) = \bar{\pi}_{k+1}^S(F)(\bar{\pi}_k^S(a))$$

iii) For each $k \in \mathbb{N}$, the range of $\overline{\pi}_k^S$ is a closed subset of $Ct^S_{\mathbb{R}}(k)$ homeomorphic to $Ct_{\mathbb{N}}(k).$

Proof. Items i) and ii) are direct consequences of Theorem 5.2.

Item iii) is proved like Corollary 3.4 with the obvious adjustment to S(0) of the first bullet point. The adjustment is obvious since each compact element of S(0) determines a closed, rational interval.

Proof of Theorem 5.2. We will adjust the construction of π_k used to prove Theorem 3.1. In order to give a sound construction we first had to define π_k and prove its properties, and then define $\bar{\pi}_k$. However, the underlying way of thinking goes the other way, we define $\bar{\pi}_k$ and then show that it is continuous by constructing an appropriate π_k .

Our definition of $\bar{\pi}_k^S$ will be almost like our definition of $\bar{\pi}_k$, the difference is that now the definition is interpreted over the intensional hierarchy while in the original case, it was interpreted over the extensional hierarchy. The challenge then is to show that $\bar{\pi}_k^S$ can be realized by an intensional object. In order to avoid repeating the details of the proof of Theorem 3.1, we will only discuss the obstacles that are new in the intensional setting.

Case k = 0: Let $\overline{\pi}_0^S(n) = n \in \mathbb{R}$ realized by $\pi_0^S(n) = n00 \cdots$. Case k = 1: Let $f : \mathbb{N} \to \mathbb{N}$ and $x \in \mathbb{R}$.

We let $\bar{\pi}_0^S(f)(x) = f(0)$ if $x \le 0$ and

$$\bar{\pi}_0^S(f)(x) = (1-y)f(n) + yf(n+1)$$

when x = n + y and $0 \le y \le 1$. Since $\mathbb{N}^{\mathbb{N}} \times \overline{S}(0)$ is zero-dimensional, it follows from Proposition 4.5 that there is a total $\pi_1^S : N(1) \times S(0) \to S(0)$ realizing $\overline{\pi}_1^S$.

Case $k \ge 2$: Let $F \in Ct_{\mathbb{N}}(k)$ and let $x \in Ct_{\mathbb{R}}^S(k-1)$.

For $\xi \in Ct_{\mathbb{N}}(k-2)$, let $f(\xi) = x(\bar{\pi}_{k-2}^{S}(\xi))$.

Let Case 1 and Case 2 correspond to the cases in the proof of Theorem 3.1. In Case 1, $f: Ct_{\mathbb{N}}(k-2) \to \mathbb{N}$ and we let

$$\pi_k^S(F)(x) = F(f).$$

In Case 2, we define

$$\bar{\pi}_{k}^{S}(F)(x) = \sum_{n=0}^{\infty} \left(d(x,n) \cdot z_{n}(x) \cdot \sum_{a \in X_{n}^{k-1}} (F(a) \cdot \mu_{n,x}(a)) \right)$$

where we use the notation from the proof of Theorem 3.1.

 $\bar{\pi}_k^S(F)(x)$ will depend continuously on the sequences $\{x(\bar{\pi}_{k-2}^S(\eta_n))\}_{n\in\mathbb{N}}$ and $\{F(a)\}_{n\in\mathbb{N},a\in X_n^{k-1}}$. Here we view η_n as an element of $Ct_{\mathbb{N}}(k-2)$ and X_n^{k-1} as a subset of $Ct_{\mathbb{N}}(k-1)$. Let \bar{G} be continuous such that

$$\bar{\pi}_k^S(F)(x) = \bar{G}(\{x(\bar{\pi}_{k-2}^S(\eta_n))\}_{n \in \mathbb{N}}, \{F(a)\}_{n \in \mathbb{N}, a \in X_n^{k-1}})$$

when we are in Case 2.

 \overline{G} is essentially of type $\mathbb{R}^{\mathbb{N}} \to \mathbb{R}$ and can be represented in the S-hierarchy by a total, continuous G. We use G to define $\pi_k^S(F)(x)$ in a continuous way from

 ${x(\tilde{\pi}_{k-2}(\eta_n))}_{n\in\mathbb{N}}$ and ${F(a)}_{n\in\mathbb{N},a\in X_n^{k-1}}$.

If we do this without any further care, we will not be able to show that the constructions from Case 1 and Case 2 match in a continuous way. However, if we use the intensional representation $F(f)00\cdots$ in Case 1, and $norm \circ G$ in Case 2 (where *norm* is the function of Lemma 5.1), we may prove continuity in the same way as we did in the proof of Theorem 3.1.

This ends our proof of the theorem.

Our next result relate the two embeddings to the core hierarchy from Normann [17]. We will not need this result elsewhere in this paper, so we assume familiarity with [17] in this proof.

Theorem 5.4. Let $x \in Ct_{\mathbb{N}}(k)$. Then $\bar{\pi}_k(x) \in Ct_{\mathbb{R}}^E(k)$ and $\bar{\pi}_k^S(x) \in Ct_{\mathbb{R}}^S(k)$ are equivalent in the sense of Normann [17].

Proof. We use induction on k, where the first two cases are trivial, so let $k \ge 2$.

Objects from these type structures are equivalent if they behave in the same way on equivalent input. We observe that $\bar{\pi}_k(F)(x)$ only depends on x restricted to the image of $\bar{\pi}_{k-2}$ and that $\bar{\pi}_k^S(F)(x)$ depends on x restricted to the image of $\bar{\pi}_{k-2}^S$ in the same way. Thus if $F \in Ct_{\mathbb{N}}(k)$ is given and $x \in Ct_{\mathbb{R}}^E$ and $y \in Ct_{\mathbb{R}}^S$ are equivalent at type k-1 we use exactly the same definition in the two cases, and the results will be two equivalent functionals.

Theorem 5.5. Assume that $Ct^{E}_{\mathbb{R}}(n+1) = Ct^{S}_{\mathbb{R}}(n+1)$. Then $\mathcal{R}(Ct_{\mathbb{N}}(n))$ is zero-dimensional.

Proof. First observe that by the assumption it follows, for purely set theoretical reasons, that $Ct^{E}_{\mathbb{R}}(k) = Ct^{S}_{\mathbb{R}}(k)$ for $k \leq n$.

Let $f: Ct_{\mathbb{N}}(n) \to \mathbb{R}$ be continuous.

Then there is a total, continuous function $h: N(n) \to E(0)$ such that

$$\rho_0^E(h(x)) = f(\rho_n^N(x))$$

for all $x \in \overline{N}(n)$.

Let Π_{n+1} be as in Lemma 3.13, and let $g = \rho_{n+1}^E(\Pi_{n+1}(h))$.

Then for all $x \in Ct_{\mathbb{N}}(n)$ we have that $f(x) = g(\bar{\pi}_n(x))$.

Since $g \in Ct^E_{\mathbb{R}}(n+1)$ it follows from the assumption that $g \in Ct^S_{\mathbb{R}}(n+1)$, which means that there is a total $\hat{g} \in \bar{S}(n+1)$ with $g = \rho_{n+1}^S(\hat{g})$.

Recall the topological space T(n) from Definition 4.12. T(n) is $\bar{S}(n)$ divided out by the consistency relation C_n on S(n).

Clearly \hat{g} will send consistent, total elements in S(n) to consistent, total elements in S(0), and consistency on $\bar{S}(0)$ is the same as identity, so \hat{q} will induce a total $\tilde{q}: T(n) \to S(0)$ $\bar{S}(0).$

Let $O \subseteq \mathbb{R}$ be open. Then

$$O^{S} = \{ y \in \overline{S}(0) ; \rho_{0}^{S}(y) \in O \}$$

is open in $\overline{S}(0)$, and since $\overline{S}(0)$ is zero-dimensional, O^S will be the union of clopen sets. Then

 $\{z \in T(n) ; \rho_0^S(\tilde{q}(z)) \in O\}$

is the union of clopen sets.

Let $\pi_n^S : N(n) \to S(n)$ be as in Theorem 5.2. We define the map π_n^T as follows: Let $x \in Ct_{\mathbb{N}}(n)$, let $y \in \overline{N}(n)$ be such that $\rho_n^N(y) = x$ and let $\pi_n^T(x) \in T(n)$ be the equivalence class of $\pi_n^S(y)$ (which is independent of the choice of y).

Then π_n^T is continuous, and

$$f^{-1}[O] = (\pi_n^T)^{-1}[\tilde{g}^{-1}[\{y \in S(0) \ ; \ \rho_0^S(y) \in O\}]].$$

It follows that $f^{-1}[O]$ is the union of clopen sets.

Since f and O was arbitrary, $\mathcal{R}(Ct_{\mathbb{N}}(n))$ will be zero-dimensional.

6. DISCUSSION

We have established the equivalence of a problem about functionals over the reals with a problem in topology, see Problem 6.2. In fact, there will be several interesting problems of topology related to the coincidense problem:

Problem 6.1. Is $Ct_{\mathbb{N}}(k)$ zero-dimensional for some or all k > 1?

Problem 6.2. Is $\mathcal{R}(Ct_{\mathbb{N}}(k))$ zero-dimensional for some or all k > 1?

Problem 6.3. Does $Ct_{\mathbb{N}}(k)$ coincide, as a topological space, with $\mathcal{R}(Ct_{\mathbb{N}}(k))$ for some or all k > 1?

Of course, if Problem 6.1 has a positive solution for a value k, then Problem 6.2 will also have a positive solution for the same k, and if both Problems 6.2 and 6.3 have positive solutions for some k, then Problem 6.1 has a positive solution for the same k.

Some effort has been put into solving Problem 6.1, but without success. At the time of writing, three possible answers to Problem 6.1 seem equally likely, the answers 'yes, 'no' and 'independent of ZFC'. There is no indication from the attempts so far that Problem 6.2 or Problem 6.3 are easier to solve. The author sees no basis for forming any conjecture.

There are topological aspects of these problems that we will not discuss in this paper, but it may well be that a deeper use of topology is what is needed to solve them. Another possibility is, as mentioned, that some or all of the problems are independent of Zermelo-Fraenkel set theory.

The statement

 $Ct_{\mathbb{N}}(2)$ is zero-dimensional

is Π_4^1 , and even, given an open set $O \subseteq Ct_{\mathbb{N}}(2)$, the statement that

O contains a nonempty clopen set

is Σ_3^1 relative to the set of compacts approximating elements in O. It is known that the truth value of some Σ_3^1 -statements can be altered by forcing. It must be stressed that so far there is nothing indicating that this will be possible in this case. The only thing we actually know is that by the Shoenfield Absoluteness Theorem, being an open subset of $Ct_{\mathbb{N}}(2)$ is absolute with respect to forcing, being a clopen subset of $Ct_{\mathbb{N}}(2)$ is absolute with respect to forcing, being a clopen subset of $Ct_{\mathbb{N}}(2)$ is absolute with respect to forcing a clopen set is the union of clopen sets, this cannot be altered by forcing. Thus forcing a counterexample must mean constructing a new open set. Open sets in $Ct_{\mathbb{N}}(2)$ constructed by a naive use of forcing, will however contain clopen sets. The use of forcing in domain theory has not been developed sufficiently far for us to comment further on this.

7. Appendix

In this appendix we will prove a special case of the approximation lemma from Normann [18]:

Theorem 7.1. Let $A \subseteq Ct_{\mathbb{N}}(k)$ and let $f : A \to \mathbb{N}$ be continuous. Then, continuously in f there are $f_n \in Ct_{\mathbb{N}}(k+1)$ for each $n \in \mathbb{N}$ such that whenever $x \in A$ and $x = \lim_{n \to \infty} x_n$ with each $x_n \in Ct_{\mathbb{N}}(k)$ we have that $f(x) = \lim_{n \to \infty} f_n(x_n)$.

Proof. Let $\bar{A} = \{x \in \bar{N}(k) ; \rho_k^N(x) \in A\}$. We will let $f \in N(k+1)$ be total on \bar{A} in the proof.

Let $\{(p_i, a_i)\}_{i \in \mathbb{N}}$ be an enumeration of all pairs (p, a) where $p \in N(k)$ is compact, p has an extension in \overline{A} and $a \in \mathbb{N}$.

If *i* and *j* are such that p_i and p_j have a joint extension in \overline{A} , let $z_{i,j}$ be one such extension. We will consider *i*, *j* as an unordered pair. In particular, $z_{i,i}$ will exist for all *i*.

Let $X_n^f = \{(p_i, a_i) ; i \leq n \text{ and for all } j \leq n, \text{ if } z_{i,j} \text{ exists, then } f(z_{i,j}) = a_i\}.$

Let $(p_i, a_i) \in Y_n^f$ if $j \leq n$ and for some $r \leq n$ we have that

- $(p_r, a_r) \in X_n^f$.
- $a_j = a_r$.
- $p_r \sqsubseteq p_j$.

It is easy to see that if $(p_j, a_j) \in Y_n^f$, $(p_{j'}, a_{j'}) \in Y_n^f$ and p_j and $p_{j'}$ are consistent, then $a_j = a_{j'}.$

By the density theorem for $Ct_{\mathbb{N}}(k+1)$ there is a total map $g = g_{Y_n^f}$ such that $g(p_j) = a_j$ whenever $(p_j, a_j) \in Y_n^f$. We let $f_n = g_{Y_n^f}$.

Now, let $x = \lim_{n \to \infty} x_n$, where $x \in \overline{A}$ and each $x_n \in \overline{N}(k)$. Let f(x) = a. Then there is an approximation p to x such that f(p) = a.

For some $r \in \mathbb{N}$ then, $(p, a) = (p_r, a_r)$. Moreover, if $r \leq n$, $i \leq n$ and $z_{r,i}$ exists, then $z_{r,i}$ extends p, so $f(z_{r,i}) = a = a_r$.

Thus $(p_r, a_r) \in X_n^f$ whenever $n \ge j$.

Let $i < r, n \ge r, (p_i, a_i) \in X_n^f$ and $a_i \ne a_r$.

Claim: a_i and x are inconsistent.

Proof. Assume that p_i and x are consistent. Then $p_i \sqcup x$ is a joint extension of p_i and p_r in A, so $z_{i,r}$ exists. Then

$$(p_i, a_i) \in X_n^f \Rightarrow f(z_{i,r}) = a_i$$

and

$$(p_r, a_r) \in X_n^f \Rightarrow f(z_{i,r}) = a_r,$$

contradicting $a_i \neq a_r$. This proves the claim.

Then there is a compact element $p \sqsubseteq x$ such that $p_r \sqsubseteq p$ and p is inconsistent with p_i whenever i < r and $a_i \neq a_r$.

For some j then, $(p, a) = (p_j, a_j)$.

If $n \ge max\{r, j\}$, it follows that $(p_j, a) \in Y_n^f$, so $f_n(p_j) = a$. Let $n_0 \ge max\{r, j\}$ be such that $n \ge n_0 \Rightarrow p_j \sqsubseteq x_n$. Then $n \ge n_0 \Rightarrow f_n(x_n) = a$, and this is what we aimed to prove.

If $y = \lim_{n \to \infty} y_n$ in $Ct_{\mathbb{N}}(k)$, there will be $x = \lim_{n \to \infty} x_n$ in $\overline{N}(k)$ such that $y = \rho_k^N(x)$ and $y_n = \rho_k^N(x_n)$ (see Proposition 2.9). Moreover, the f_n constructed only depends on f restricted to \overline{A} , and is thus definable from $\rho_A(f) : A \to \mathbb{N}$, where

$$\rho_A(f)(\rho_k^N(x)) = f(x)$$

Actually, f_n is definable from f restricted to

$$\{\rho_k^N(z_{i,j}) ; i \leq n, j \leq n \text{ and } z_{i,j} \text{ exists } \},\$$

so f_n depends continuously on f. Now, by Proposition 2.10 , every continuous function from A to \mathbb{N} will be of the form $\rho_A(f)$, so we are through.

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